

INTERPRETABILITY

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Definition 0.1 (Interpretation). *Let L and K be languages. Let \mathcal{M} be an L -structure and \mathcal{N} be a K -structure. If \mathcal{M}^* is an L^* -structure, where $K \subset L^*$ such that there is a map $\cdot^* : L \rightarrow L^*$ such that $\mathcal{M} \models \varphi \iff \mathcal{M}^* \models \varphi^*$ then \mathcal{M}^* and every structure isomorphic to it is an **interpretation** of \mathcal{M} in \mathcal{N} .*

Remark 0.1. *By the interpretation lemma, to define the map \cdot^* it suffices to show that every constant, function and relation of L is definable in K .*

Example 0.2. *Interpreting $\mathbb{C} = (\mathbb{C}, 0, 1, +, -, \cdot)$ in \mathbb{R} as a $(\mathbb{R}, 0, 1, +, -, \cdot)$ structure.*

Given $x \in \mathbb{R}$ denote by x^{-1} the unique real number y such that $xy = 1$. Note that addition $(+(z, u, v))$, 0, multiplication $(\cdot(z, u, v))$, 1 and inverses are definable, in particular, if $0^* = (0, 0)$ and $1^* = (1, 0)$ then

$$\begin{aligned} +(z, u, v) \xrightarrow{*} \varsigma(x, y, a, b, c, d) &\iff (x = a + c) \wedge (y = b + d) \\ \cdot(z, u, v) \xrightarrow{*} \mu(x, y, a, b, c, d) &\iff (x = ac - bd) \wedge (y = ad + bc) \\ \exists^{-1} z'(zz' = 1) \xrightarrow{*} \iota(x, y, a, b) &\iff \exists^{-1} xy(x = a(aa + bb)^{-1}) \wedge (y = -b(aa + bb)^{-1}). \end{aligned}$$

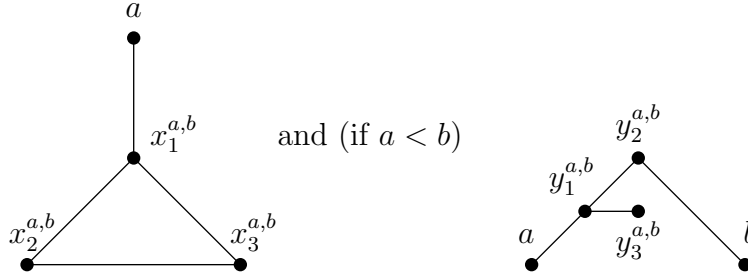
We therefore may use the following abbreviations for these operations, understanding (a, b) to be $a + bi$:

$$\begin{aligned} (a, b) + (c, d) &= (a + c, b + d) \\ (a, b)(c, d) &= (ac - bd, ad + bc) \\ (a, b)^{-1} &= (a(aa + bb)^{-1}, -b(aa + bb)^{-1}). \end{aligned}$$

Example 0.3. *Interpreting linear orders in graphs as $\{R\}$ -structures.*

Let $\triangleleft(x, y, u, v, w)$ be the formula asserting that x, y, u, v and w are distinct, there are edges $(x, u), (u, v), (u, w)$ and (v, y) and these are the only edges involving vertices u, v and w . Define $\triangleleft(x, y, u, v)$ in a similar fashion. Now define $\triangleleft(z)$ as $\exists uvw \triangleleft(z, u, v, w)$ and $\triangleleft(z), \triangleleft(z), \triangleleft(z), \triangleleft(z), \triangleleft(z)$ analogously. Finally, let $S = \{ \triangleleft(z), \triangleleft(z), \triangleleft(z), \triangleleft(z), \triangleleft(z), \triangleleft(z) \}$.

If $(A, <)$ is a linear order then define a graph G_A such that for each $(a, b) \in A^2$



are subgraphs of G_A . Further construct G_A so that these are the only edges which involve any x_i^a or $y_i^{a,b}$.

Lemma 0.4. *If $(A, <)$ is a linear order then for all vertices $x \in G_A$ there is exactly one $\varphi(x) \in S$ such that $G_A \models \varphi(x)$.*

Let T be the theory in the language of graphs with the following axioms:

- (1) R is symmetric and irreflexive
- (2) Exactly one $\varphi \in S$ is true for every x .
- (3) If $\begin{array}{c} \circ \\ \diagup \quad \diagdown \\ x \end{array}$ and $\begin{array}{c} \circ \\ \diagup \quad \diagdown \\ y \end{array}$ then $\neg R(x, y)$.
- (4) If $\exists uvw \begin{array}{c} \diagdown \quad \diagup \\ x, y, u, v, w \end{array}$ then $\forall u_1 v_1 w_1 \neg \begin{array}{c} \diagdown \quad \diagup \\ y, x, u_1, v_1, w_1 \end{array}$.
- (5) If $\exists uvw \begin{array}{c} \diagdown \quad \diagup \\ x, y, u, v, w \end{array}$ and $\exists u_1 v_1 w_1 \begin{array}{c} \diagdown \quad \diagup \\ y, z, u_1, v_1, w_1 \end{array}$ then $\exists u_2 v_2 w_2 \begin{array}{c} \diagdown \quad \diagup \\ x, z, u_2, v_2, w_2 \end{array}$.
- (6) If $\begin{array}{c} \circ \\ \diagup \quad \diagdown \\ x \end{array}$ and $\begin{array}{c} \circ \\ \diagup \quad \diagdown \\ y \end{array}$ then either $x = y$,
 $\exists uvw \begin{array}{c} \diagdown \quad \diagup \\ x, y, u, v, w \end{array}$ or $\exists uvw \begin{array}{c} \diagdown \quad \diagup \\ y, x, u, v, w \end{array}$
- (7) If $\begin{array}{c} \diagdown \quad \diagup \\ x, u, v, w \end{array} \wedge \begin{array}{c} \diagdown \quad \diagup \\ x, u_1, v_1, w_1 \end{array}$ then $u = u_1$ and $v = v_1 \wedge w = w_1$ or $v = w_1 \wedge w = v_1$.
- (8) If $\begin{array}{c} \diagdown \quad \diagup \\ x, y, u, v, w \end{array} \wedge \begin{array}{c} \diagdown \quad \diagup \\ x, y, u_1, v_1, w_1 \end{array}$ then $u = u_1, v = v_1$ and $w = w_1$.

Remark 0.5. *If $(A, <)$ is a linear order then $G_A \models T$.*

Remark 0.6. *We may also interpret other common order-related axioms in this language. Using the interpretation map above we see that density $(\forall xy(x < y) \rightarrow \exists z(x < z \wedge z < y))$ becomes:*

$$\begin{aligned} \forall xy \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ x \end{array} \wedge \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ y \end{array} \wedge \exists uvw \begin{array}{c} \diagdown \quad \diagup \\ x, y, u, v, w \end{array} \rightarrow \\ \exists z \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ z \end{array} \wedge \\ \exists u_x v_x w_x \begin{array}{c} \diagdown \quad \diagup \\ u_x \end{array} \wedge \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ v_x \end{array} \wedge \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ w_x \end{array} \wedge \begin{array}{c} \diagdown \quad \diagup \\ x, z, u_x, v_x, w_x \end{array} \wedge \\ \exists u_y v_y w_y \begin{array}{c} \diagdown \quad \diagup \\ u_y \end{array} \wedge \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ v_y \end{array} \wedge \begin{array}{c} \circ \\ \diagup \quad \diagdown \\ w_y \end{array} \wedge \begin{array}{c} \diagdown \quad \diagup \\ z, y, u_y, v_y, w_y \end{array}. \end{aligned}$$